

Majorization Minimization (MM) and Block Coordinate Descent (BCD)

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Outline

- Majorization Minimization (MM)
 - Convergence
 - Applications
- Block Coordinate Descent (BCD)
 - Applications
 - Convergence
- Summary

Majorization Minimization

Consider the following problem

$$\min_{\mathbf{x}} f(\mathbf{x}) \quad \text{s.t. } \mathbf{x} \in \mathcal{X} \quad (1)$$

where \mathcal{X} is a closed convex set; $f(\cdot)$ may be non-convex and/or nonsmooth.

- **Challenge:** For a general $f(\cdot)$, problem (1) can be difficult to solve.
- **Majorization Minimization:** Iteratively generate $\{\mathbf{x}^r\}$ as follows

$$\mathbf{x}^r \in \min_{\mathbf{x}} u(\mathbf{x}, \mathbf{x}^{r-1}) \quad \text{s.t. } \mathbf{x} \in \mathcal{X} \quad (2)$$

where $u(\mathbf{x}, \mathbf{x}^{r-1})$ is a surrogate function of $f(\mathbf{x})$, satisfying

1. $u(\mathbf{x}, \mathbf{x}^r) \geq f(\mathbf{x}), \quad \forall \mathbf{x}^r, \mathbf{x} \in \mathcal{X};$
2. $u(\mathbf{x}^r, \mathbf{x}^r) = f(\mathbf{x}^r);$

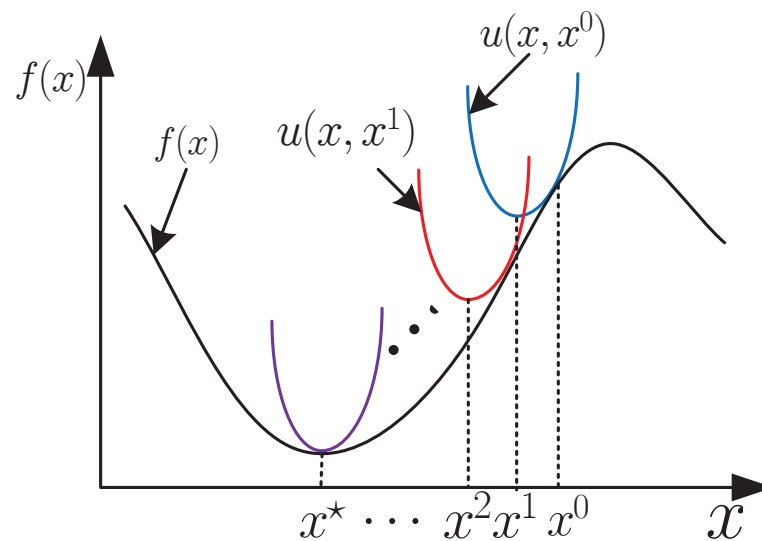


Figure 1: An pictorial illustration of MM algorithm.

Property 1. $\{f(\mathbf{x}^r)\}$ is nonincreasing, i.e., $f(\mathbf{x}^r) \leq f(\mathbf{x}^{r-1})$, $\forall r = 1, 2, \dots$

Proof. $f(\mathbf{x}^r) \leq u(\mathbf{x}^r, \mathbf{x}^{r-1}) \leq u(\mathbf{x}^{r-1}, \mathbf{x}^{r-1}) = f(\mathbf{x}^{r-1})$ ■

- The nonincreasing property of $\{f(\mathbf{x}^r)\}$ implies that $f(\mathbf{x}^r) \rightarrow \bar{f}$. But how about the convergence of the iterates $\{\mathbf{x}^r\}$?

Technical Preliminaries

- **Limit point:** \bar{x} is a limit point of $\{x_k\}$ if there exists a subsequence of $\{x_k\}$ that converges to \bar{x} . Note that every bounded sequence in \mathbb{R}^n has a limit point (or convergent subsequence);
- **Directional derivative:** Let $f : \mathcal{D} \rightarrow \mathbb{R}$ be a function where $\mathcal{D} \subseteq \mathbb{R}^m$ is a convex set. The directional derivative of f at point x in direction \mathbf{d} is defined by

$$f'(\mathbf{x}; \mathbf{d}) \triangleq \liminf_{\lambda \downarrow 0} \frac{f(\mathbf{x} + \lambda \mathbf{d}) - f(\mathbf{x})}{\lambda}.$$

– If f is differentiable, then $f'(\mathbf{x}; \mathbf{d}) = \mathbf{d}^T \nabla f(\mathbf{x})$.

- **Stationary point:** $x \in \mathcal{X}$ is a stationary point of $f(\cdot)$ if

$$f'(\mathbf{x}; \mathbf{d}) \geq 0, \quad \forall \mathbf{d} \text{ such that } \mathbf{x} + \mathbf{d} \in \mathcal{D}. \quad (3)$$

- A stationary point may be a local min., a local max. or a saddle point;
- If $\mathcal{D} = \mathbb{R}^n$ and f is differentiable, then (3) $\iff \nabla f(\mathbf{x}) = \mathbf{0}$.

Convergence of MM

- **Assumption 1** $u(\cdot, \cdot)$ satisfies the following conditions

$$\left\{ \begin{array}{l} u(\mathbf{y}, \mathbf{y}) = f(\mathbf{y}), \quad \forall \mathbf{y} \in \mathcal{X}, \\ u(\mathbf{x}, \mathbf{y}) \geq f(\mathbf{x}), \quad \forall \mathbf{x}, \mathbf{y} \in \mathcal{X}, \\ u'(\mathbf{x}, \mathbf{y}; \mathbf{d})|_{\mathbf{x}=\mathbf{y}} = f'(\mathbf{y}; \mathbf{d}), \quad \forall \mathbf{d} \text{ with } \mathbf{y} + \mathbf{d} \in \mathcal{X}, \\ u(\mathbf{x}, \mathbf{y}) \text{ is continuous in } (\mathbf{x}, \mathbf{y}), \end{array} \right. \begin{array}{l} (4a) \\ (4b) \\ (4c) \\ (4d) \end{array}$$

- (4c) means the 1st order local behavior of $u(\cdot, \mathbf{x}^{r-1})$ is the same as $f(\cdot)$.

Convergence of MM

Theorem 1. [Razaviyayn-Hong-Luo] Assume that Assumption 1 is satisfied. Then every limit point of the iterates generated by MM algorithm is a stationary point of problem (1).

Proof. From **Property 1**, we know that $f(\mathbf{x}^{r+1}) \leq u(\mathbf{x}^{r+1}, \mathbf{x}^r) \leq u(\mathbf{x}, \mathbf{x}^r)$, $\forall \mathbf{x} \in \mathcal{X}$. Now assume that there exists a subsequence $\{\mathbf{x}^{r_j}\}$ of $\{\mathbf{x}^r\}$ converging to a limit point \mathbf{z} , i.e., $\lim_{j \rightarrow \infty} \mathbf{x}^{r_j} = \mathbf{z}$. Then

$$u(\mathbf{x}^{r_{j+1}}, \mathbf{x}^{r_{j+1}}) = f(\mathbf{x}^{r_{j+1}}) \leq f(\mathbf{x}^{r_{j+1}}) \leq u(\mathbf{x}^{r_{j+1}}, \mathbf{x}^{r_j}) \leq u(\mathbf{x}, \mathbf{x}^{r_j}), \quad \forall \mathbf{x} \in \mathcal{X}.$$

Letting $j \rightarrow \infty$, we obtain $u(\mathbf{z}, \mathbf{z}) \leq u(\mathbf{x}, \mathbf{z})$, $\forall \mathbf{x} \in \mathcal{X}$, which implies that

$$u'(\mathbf{x}, \mathbf{z}; \mathbf{d})|_{\mathbf{x}=\mathbf{z}} \geq 0, \quad \forall \mathbf{z} + \mathbf{d} \in \mathcal{X}.$$

Combining the above inequality with (4c) (i.e., $u'(\mathbf{x}, \mathbf{y}; \mathbf{d})|_{\mathbf{x}=\mathbf{y}} = f'(\mathbf{y}; \mathbf{d})$, $\forall \mathbf{d}$ with $\mathbf{y} + \mathbf{d} \in \mathcal{X}$), we have

$$f'(\mathbf{z}; \mathbf{d}) \geq 0, \quad \forall \mathbf{z} + \mathbf{d} \in \mathcal{X}.$$

Applications — Nonnegative Least Squares

In many engineering applications, we encounter the following problem

$$\text{(NLS)} \quad \min_{\mathbf{x} \geq \mathbf{0}} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2 \quad (5)$$

where $\mathbf{b} \in \mathbb{R}_+^m$, $\mathbf{b} \neq \mathbf{0}$, and $\mathbf{A} \in \mathbb{R}_{++}^{m \times n}$.

- It's an LS problem with nonnegative constraints, so the conventional LS solution may not be feasible for (5).
- A simple multiplicative updating algorithm:

$$\mathbf{x}_l^r = c_l^r \mathbf{x}_l^{r-1}, \quad l = 1, \dots, n \quad (6)$$

where \mathbf{x}_l^r is the l th component of \mathbf{x}^r , and $c_l^r = \frac{[\mathbf{A}^T \mathbf{b}]_l}{[\mathbf{A}^T \mathbf{A} \mathbf{x}^{r-1}]_l}$.

- Starting with $\mathbf{x}^0 > \mathbf{0}$, then all \mathbf{x}^r generated by (6) are nonnegative.

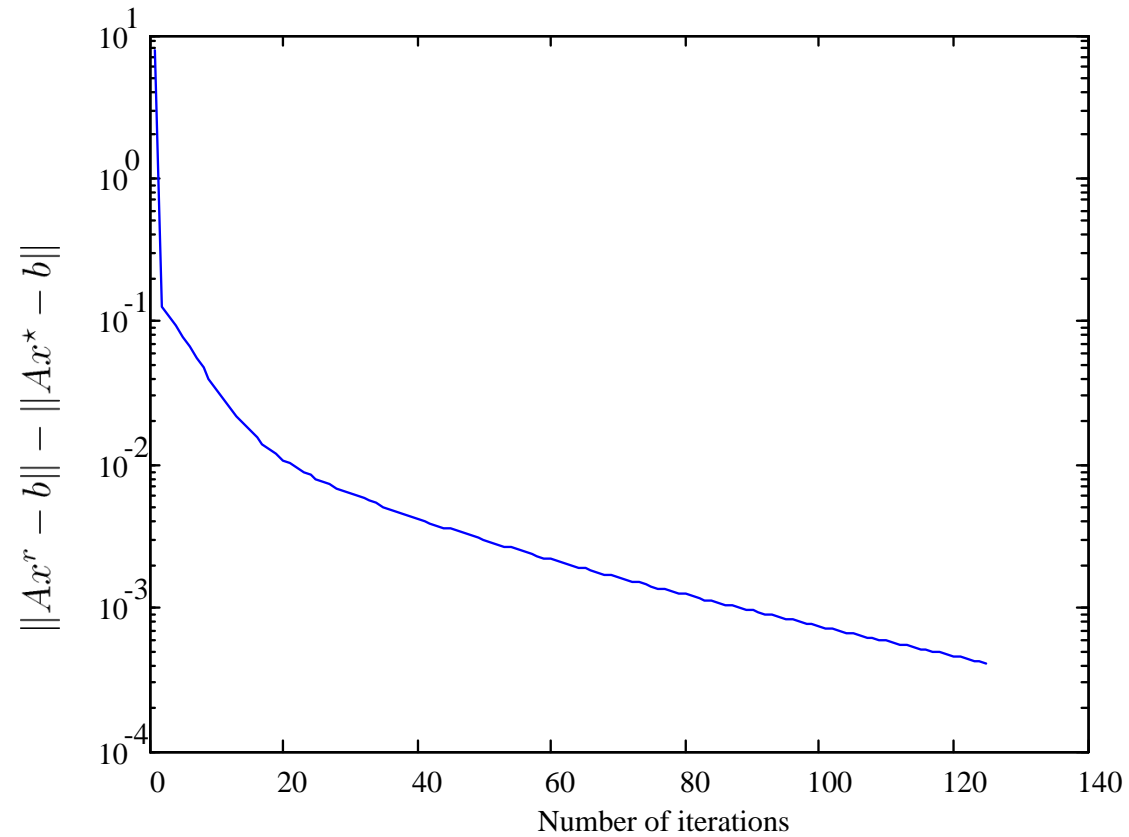


Figure 2: $\|\mathbf{A}\mathbf{x}^r - \mathbf{b}\|_2$ vs. the number of iterations.

- Usually the multiplicative update converges within a few tens of iterations.

- **MM interpretation:** Let $f(\mathbf{x}) \triangleq \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2$. The multiplicative update essentially solves the following problem

$$\min_{\mathbf{x} \geq \mathbf{0}} u(\mathbf{x}, \mathbf{x}^{r-1})$$

where

$$u(\mathbf{x}, \mathbf{x}^{r-1}) \triangleq f(\mathbf{x}^{r-1}) + (\mathbf{x} - \mathbf{x}^{r-1})^T \nabla f(\mathbf{x}^{r-1}) + \frac{1}{2} (\mathbf{x} - \mathbf{x}^{r-1})^T \Phi(\mathbf{x}^{r-1}) (\mathbf{x} - \mathbf{x}^{r-1}),$$

$$\Phi(\mathbf{x}^{r-1}) = \text{Diag} \left(\frac{[\mathbf{A}^T \mathbf{A} \mathbf{x}^{r-1}]_1}{x_1^{r-1}}, \dots, \frac{[\mathbf{A}^T \mathbf{A} \mathbf{x}^{r-1}]_n}{x_n^{r-1}} \right).$$

– Observations:

$$\begin{cases} u(\mathbf{x}, \mathbf{x}^{r-1}) \text{ is quadratic approx. of } f(\mathbf{x}), \\ \Phi(\mathbf{x}^{r-1}) \succeq \mathbf{A}^T \mathbf{A}, \end{cases} \implies \begin{cases} u(\mathbf{x}, \mathbf{x}^{r-1}) \geq f(\mathbf{x}), \quad \forall \mathbf{x} \in \mathbb{R}^n, \\ u(\mathbf{x}^{r-1}, \mathbf{x}^{r-1}) = f(\mathbf{x}^{r-1}). \end{cases}$$

- The multiplicative update converges to an optimal solution of NLS (by the MM convergence in **Theorem 1** and convexity of NLS).

Applications — Convex-Concave Procedure/ DC Programming

- Suppose that $f(\mathbf{x})$ has the following form

$$f(\mathbf{x}) = g(\mathbf{x}) - h(\mathbf{x}),$$

where $g(\mathbf{x})$ and $h(\mathbf{x})$ are convex and differentiable. Thus, $f(\mathbf{x})$ is in general nonconvex.

- **DC Programming:** Construct $u(\cdot, \cdot)$ as

$$u(\mathbf{x}, \mathbf{x}^r) = g(\mathbf{x}) - \underbrace{\left(h(\mathbf{x}^r) + \nabla_{\mathbf{x}} h(\mathbf{x}^r)^T (\mathbf{x} - \mathbf{x}^r) \right)}_{\text{linearization of } h \text{ at } \mathbf{x}^r}$$

- By the 1st order condition of $h(\mathbf{x})$, it's easy to show that

$$u(\mathbf{x}, \mathbf{x}^r) \geq f(\mathbf{x}), \quad \forall \mathbf{x} \in \mathcal{X}, \quad u(\mathbf{x}^r, \mathbf{x}^r) = f(\mathbf{x}^r).$$

- **Sparse Signal Recovery by DC Programming**

$$\min_{\mathbf{x}} \|\mathbf{x}\|_0 \quad \text{s.t. } \mathbf{y} = \mathbf{A}\mathbf{x} \quad (7)$$

– Apart from the popular ℓ_1 approximation, consider the following concave approximation:

$$\min_{\mathbf{x} \in \mathbb{R}^n} \sum_{i=1}^n \log(1 + |x_i|/\epsilon) \quad \text{s.t. } \mathbf{y} = \mathbf{A}\mathbf{x},$$

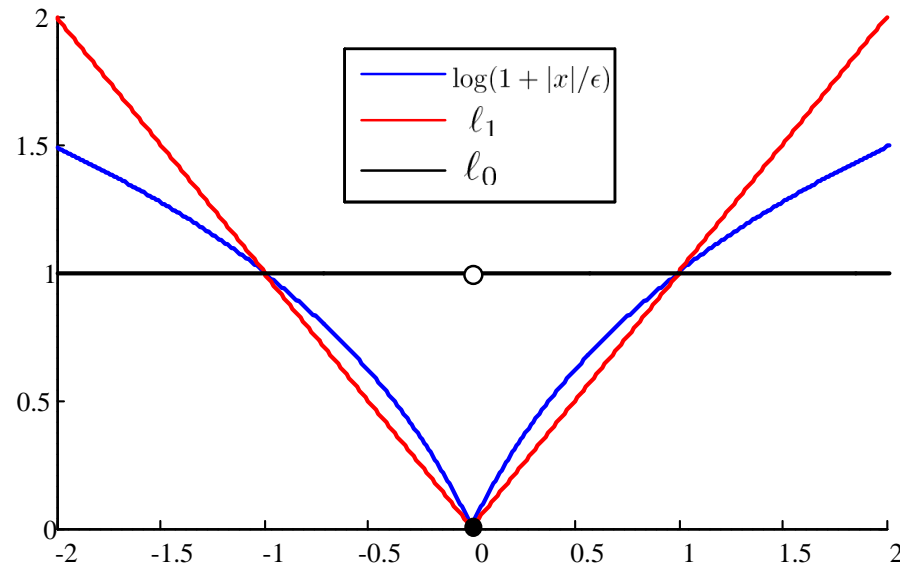


Figure 3: $\log(1 + |x|/\epsilon)$ promotes more sparsity than ℓ_1

$$\min_{\mathbf{x} \in \mathbb{R}^n} \sum_{i=1}^n \log(1 + |x_i|/\epsilon) \quad \text{s.t. } \mathbf{y} = \mathbf{A}\mathbf{x},$$

which can be equivalently written as

$$\min_{\mathbf{x}, \mathbf{z} \in \mathbb{R}^n} \sum_{i=1}^n \log(z_i + \epsilon) \quad \text{s.t. } \mathbf{y} = \mathbf{A}\mathbf{x}, \quad |x_i| \leq z_i, \quad i = 1, \dots, n \quad (8)$$

- Problem (8) minimizes a concave objective, so it's a special case of DC programming ($g(\mathbf{x}) = 0$). Linearizing the concave function at $(\mathbf{x}^r, \mathbf{z}^r)$ yields

$$(\mathbf{x}^{r+1}, \mathbf{z}^{r+1}) = \arg \min \sum_{i=1}^n \frac{z_i}{z_i^r + \epsilon} \quad \text{s.t. } \mathbf{y} = \mathbf{A}\mathbf{x}, \quad |x_i| \leq z_i, \quad i = 1, \dots, n$$

- We solve a sequence of reweighted ℓ_1 problems.

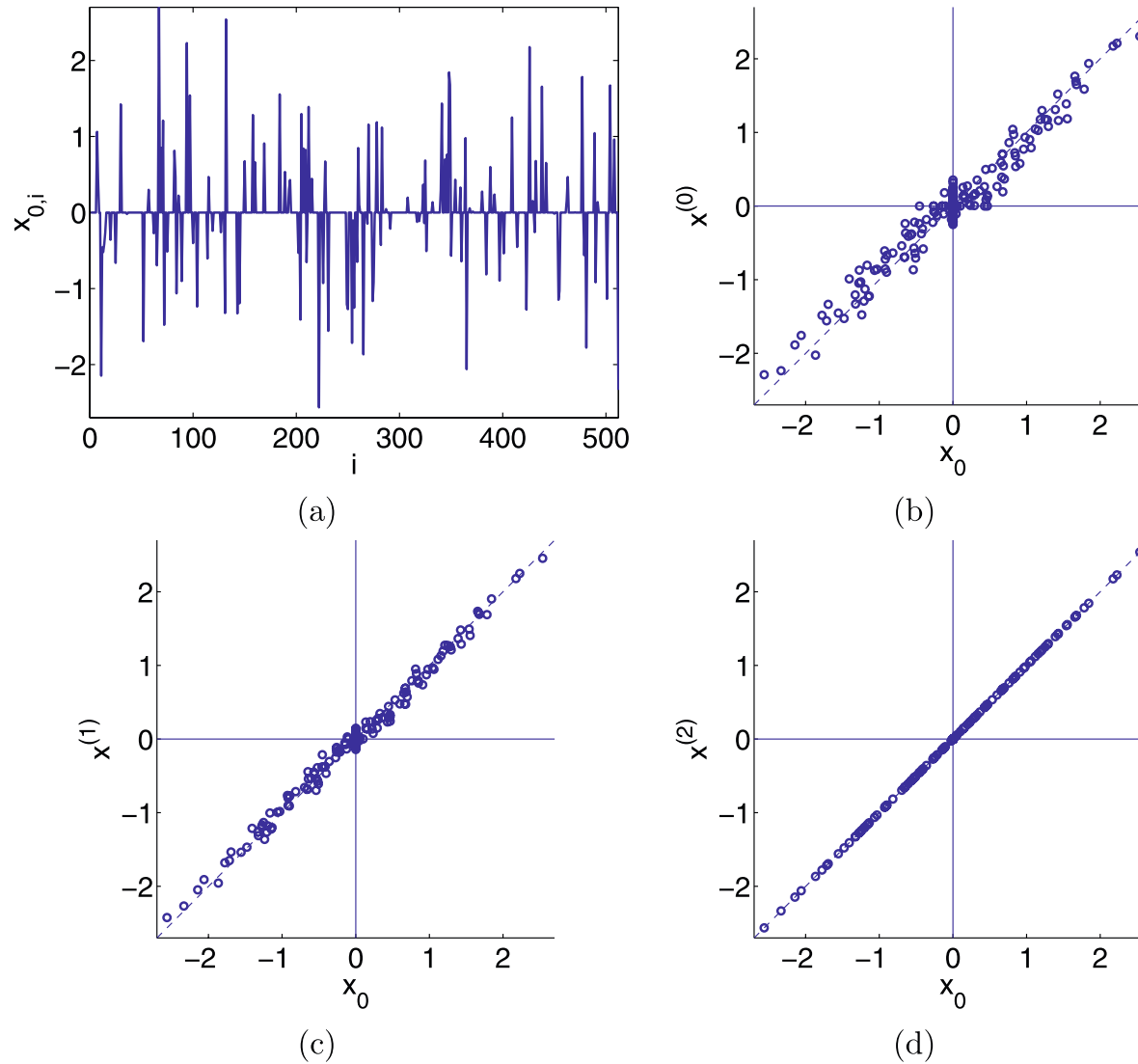


Fig. 2 Sparse signal recovery through reweighted ℓ_1 iterations. **(a)** Original length $n = 512$ signal x_0 with 130 spikes. **(b)** Scatter plot, coefficient-by-coefficient, of x_0 versus its reconstruction $x^{(0)}$ using unweighted ℓ_1 minimization. **(c)** Reconstruction $x^{(1)}$ after the first reweighted iteration. **(d)** Reconstruction $x^{(2)}$ after the second reweighted iteration

Applications — $\ell_2 - \ell_p$ Optimization

- Many problems involve solving the following problem (e.g., basis-pursuit denoising)

$$\min_{\mathbf{x}} f(\mathbf{x}) \triangleq \frac{1}{2} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 + \mu \|\mathbf{x}\|_p \quad (9)$$

where $p \geq 1$.

- If $\mathbf{A} = \mathbf{I}$ or \mathbf{A} is unitary, optimal \mathbf{x}^* is computed in closed-form as

$$\mathbf{x}^* = \mathbf{A}^T \mathbf{y} - \text{Proj}_C(\mathbf{A}^T \mathbf{y})$$

where $C \triangleq \{\mathbf{x} : \|\mathbf{x}\|_{p^*} \leq \mu\}$, $\|\cdot\|_{p^*}$ is the dual norm of $\|\cdot\|_p$ and Proj_C denotes the projection operator. In particular, for $p = 1$

$$x_i^* = \text{soft}(y_i, \mu), \quad i = 1, \dots, n$$

where $\text{soft}(u, a) \triangleq \text{sign}(u) \max\{|u| - a, 0\}$ denotes a *soft-thresholding* operation.

- For general \mathbf{A} , there is no simple closed-form solution for (9).

- **MM for $\ell_2 - \ell_p$ Problem:** Consider a modified $\ell_2 - \ell_p$ problem

$$\min_{\mathbf{x}} u(\mathbf{x}, \mathbf{x}^r) \triangleq f(\mathbf{x}) + \text{dist}(\mathbf{x}, \mathbf{x}^r) \quad (10)$$

where $\text{dist}(\mathbf{x}, \mathbf{x}^r) \triangleq \frac{c}{2} \|\mathbf{x} - \mathbf{x}^r\|_2^2 - \frac{1}{2} \|\mathbf{A}\mathbf{x} - \mathbf{A}\mathbf{x}^r\|_2^2$ and $c > \lambda_{\max}(\mathbf{A}^T \mathbf{A})$.

- $\text{dist}(\mathbf{x}, \mathbf{x}^r) \geq 0 \quad \forall \mathbf{x} \implies u(\mathbf{x}, \mathbf{x}^r)$ majorizes $f(\mathbf{x})$.
- $u(\mathbf{x}, \mathbf{x}^r)$ can be reexpressed as

$$u(\mathbf{x}, \mathbf{x}^r) = \frac{c}{2} \|\mathbf{x} - \bar{\mathbf{x}}^r\|_2^2 + \mu \|\mathbf{x}\|_p + \text{const.},$$

where

$$\bar{\mathbf{x}}^r = \frac{1}{c} \mathbf{A}^T (\mathbf{y} - \mathbf{A}\mathbf{x}^r) + \mathbf{x}^r.$$

- The modified $\ell_2 - \ell_p$ problem (10) has a simple soft-thresholding solution.
- Repeatedly solving problem (10) leads to an optimal solution of the $\ell_2 - \ell_p$ problem (by the MM convergence in **Theorem 1**)

Applications — Expectation Maximization (EM)

- Consider an ML estimate of θ , given the random observation w

$$\hat{\theta}_{\text{ML}} = \arg \min_{\theta} -\ln p(w|\theta)$$

- Suppose that there are some missing data or hidden variables z in the model. Then, EM algorithm iteratively compute an ML estimate $\hat{\theta}$ as follows:

– E-step:

$$g(\theta, \theta^r) \triangleq \mathbb{E}_{z|w, \theta^r} \{\ln p(w, z|\theta)\}$$

– M-step:

$$\theta^{r+1} = \arg \max_{\theta} g(\theta, \theta^r)$$

– repeat the above two steps until convergence.

- EM algorithm generates a nonincreasing sequence of $\{-\ln p(w|\theta^r)\}$.
- EM algorithm can be interpreted by MM.

- MM interpretation of EM algorithm:

$$\begin{aligned}
& -\ln p(w|\theta) \\
&= -\ln \mathbb{E}_{z|\theta} p(w|z, \theta) \\
&= -\ln \mathbb{E}_{z|\theta} \left[\frac{p(z|w, \theta^r) p(w|z, \theta)}{p(z|w, \theta^r)} \right] \\
&= -\ln \mathbb{E}_{z|w, \theta^r} \left[\frac{p(z|\theta) p(w|z, \theta)}{p(z|w, \theta^r)} \right] \quad (\text{interchange the integrations}) \\
&\leq -\mathbb{E}_{z|w, \theta^r} \ln \left[\frac{p(z|\theta) p(w|z, \theta)}{p(z|w, \theta^r)} \right] \quad (\text{Jensen's inequality}) \\
&= -\mathbb{E}_{z|w, \theta^r} \ln p(w, z|\theta) + \mathbb{E}_{z|w, \theta^r} \ln p(z|w, \theta^r) \tag{11a} \\
&\triangleq u(\theta, \theta^r)
\end{aligned}$$

- $u(\theta, \theta^r)$ majorizes $-\ln p(w|\theta)$, and $-\ln p(w|\theta^r) = u(\theta^r, \theta^r)$;
- E-step essentially constructs $u(\theta, \theta^r)$;
- M-step minimizes $u(\theta, \theta^r)$ (note θ appears in the 1st term of (11a) only).

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Block Coordinate Descent

- Consider the following problem

$$\min_{\mathbf{x}} f(\mathbf{x}) \quad \text{s.t. } \mathbf{x} \in \mathcal{X} = \mathcal{X}_1 \times \mathcal{X}_2 \times \dots \times \mathcal{X}_m \subseteq \mathbb{R}^n \quad (12)$$

where each $\mathcal{X}_i \subseteq \mathbb{R}^{n_i}$ is closed, nonempty and convex.

- **BCD Algorithm:**

- 1: Find a feasible point $\mathbf{x}^0 \in \mathcal{X}$ and set $r = 0$
- 2: **repeat**
- 3: $r = r + 1, i = (r - 1 \bmod m) + 1$
- 4: Let $\mathbf{x}_i^* \in \arg \min_{\mathbf{x} \in \mathcal{X}_i} f(\mathbf{x}_1^{r-1}, \dots, \mathbf{x}_{i-1}^{r-1}, \mathbf{x}, \mathbf{x}_{i+1}^{r-1}, \dots, \mathbf{x}_m^{r-1})$
- 5: Set $\mathbf{x}_i^r = \mathbf{x}_i^*$ and $\mathbf{x}_k^r = \mathbf{x}_k^{r-1}, \forall k \neq i$
- 6: **until** some convergence criterion is met

- Merits of BCD

1. each subproblem is much easier to solve, or even has a closed-form solution;
2. The objective value is nonincreasing along the BCD updates;
3. it allows parallel or distributed implementations.

Applications — $\ell_2 - \ell_1$ Optimization Problem

- Let us revisit the $\ell_2 - \ell_1$ problem

$$\min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}) \triangleq \frac{1}{2} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 + \mu \|\mathbf{x}\|_1 \quad (13)$$

- Apart from MM, BCD is another efficient approach to solve (13):

- Optimize x_k while fixing $x_j = x_j^r, \forall j \neq k$:

$$\min_{x_k} f_k(x_k) \triangleq \frac{1}{2} \left\| \mathbf{y} - \underbrace{\sum_{j \neq k} \mathbf{a}_j x_j^r}_{\triangleq \bar{\mathbf{y}}} - \mathbf{a}_k x_k \right\|_2^2 + \mu |x_k|$$

- The optimal x_k has a closed form:

$$x_k^* = \text{soft} \left(\mathbf{a}_k^T \bar{\mathbf{y}} / \|\mathbf{a}_k\|^2, \mu / \|\mathbf{a}_k\|^2 \right)$$

- Cyclically update $x_k, k = 1, \dots, n$ until convergence.

Applications — Iterative Water-filling for MIMO MAC Sum Capacity Maximization

- **MIMO Channel Capacity Maximization**

- MIMO received signal model:

$$\mathbf{y}(t) = \mathbf{H}\mathbf{x}(t) + \mathbf{n}(t)$$

where

$$\mathbf{x}(t) \in \mathbb{C}^N$$

Tx signal

$$\mathbf{H} \in \mathbb{C}^{N \times N}$$

MIMO channel matrix

$$\mathbf{n}(t) \in \mathbb{C}^N$$

standard additive Gaussian noise, i.e., $\mathbf{n}(t) \sim \mathcal{CN}(\mathbf{0}, \mathbf{I})$.



Figure 4: MIMO system model.

- MIMO channel capacity:

$$C(\mathbf{Q}) = \log \det (\mathbf{I} + \mathbf{H}\mathbf{Q}\mathbf{H}^H)$$

where $\mathbf{Q} = \mathbb{E}\{\mathbf{x}(t)\mathbf{x}(t)^H\}$ is the covariance of the tx signal.

- MIMO channel capacity maximization:

$$\max_{\mathbf{Q} \succeq \mathbf{0}} \log \det (\mathbf{I} + \mathbf{H}\mathbf{Q}\mathbf{H}^H) \quad \text{s.t. } \text{Tr}(\mathbf{Q}) \leq P$$

where $P > 0$ is the transmit power budget.

- The optimal \mathbf{Q}^* is given by the well-known *water-filling* solution, i.e.,

$$\mathbf{Q}^* = \mathbf{V}\text{Diag}(\mathbf{p}^*)\mathbf{V}^H$$

where $\mathbf{H} = \mathbf{U}\text{Diag}(\sigma_1, \dots, \sigma_N)\mathbf{V}^H$ is the SVD of \mathbf{H} , and $\mathbf{p}^* = [p_1^*, \dots, p_N^*]$ is the power allocation with $p_i^* = \max(0, \mu - 1/\sigma_i^2)$ and $\mu \geq 0$ being the water-level such that $\sum_i p_i^* = P$.

- **MIMO Multiple-Access Channel (MAC) Sum-Capacity Maximization**

- Multiple transmitters simultaneously communicate with one receiver:

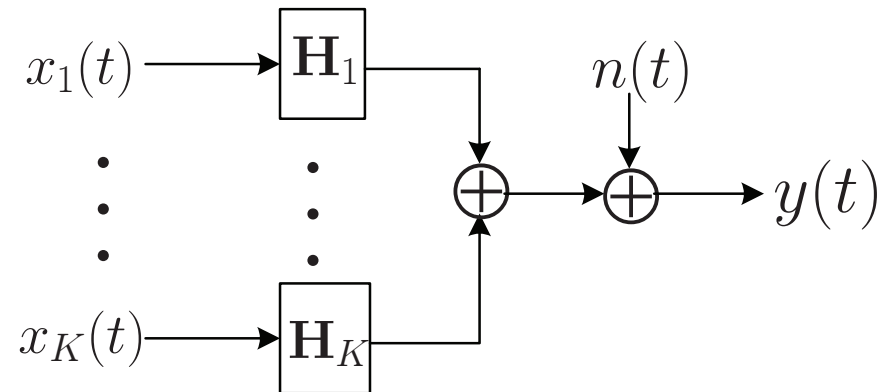


Figure 5: MIMO multiple-access channel (MAC).

- Received signal model:

$$\mathbf{y}(t) = \sum_{k=1}^K \mathbf{H}_k \mathbf{x}_k(t) + \mathbf{n}(t)$$

- MAC sum capacity:

$$C_{\text{MAC}}(\{\mathbf{Q}_k\}_{k=1}^K) = \log \det \left(\sum_{k=1}^K \mathbf{H}_k \mathbf{Q}_k \mathbf{H}_k^H + \mathbf{I} \right)$$

- MAC sum capacity maximization:

$$\begin{aligned} \max_{\{\mathbf{Q}_k\}_{k=1}^K} \quad & \log \det \left(\sum_{k=1}^K \mathbf{H}_k \mathbf{Q}_k \mathbf{H}_k^H + \mathbf{I} \right) \\ \text{s.t.} \quad & \text{Tr}(\mathbf{Q}_k) \leq P_k, \quad \mathbf{Q}_k \succeq \mathbf{0}, \quad k = 1, \dots, K \end{aligned} \quad (14)$$

- Problem (14) is convex w.r.t. $\{\mathbf{Q}_k\}$, but it has no simple closed-form solution.
- Alternatively, we can apply BCD to (14) and cyclically update \mathbf{Q}_k while fixing \mathbf{Q}_j for $j \neq k$

$$\begin{aligned} (\triangle) \quad & \max_{\mathbf{Q}_k} \log \det \left(\mathbf{H}_k \mathbf{Q}_k \mathbf{H}_k^H + \Phi \right) \\ \text{s.t.} \quad & \text{Tr}(\mathbf{Q}_k) \leq P_k, \quad \mathbf{Q}_k \succeq \mathbf{0}, \end{aligned}$$

where $\Phi = \sum_{j \neq k} \mathbf{H}_j \mathbf{Q}_j \mathbf{H}_j^H + \mathbf{I}$

- (\triangle) has a closed-form water-filling solution, just like the previous single-user MIMO case.

- An alternative low-rank matrix completion formulation [Wen-Yin-Zhang]:

$$(\triangle) \quad \min_{\mathbf{X}, \mathbf{Y}, \mathbf{Z}} \frac{1}{2} \|\mathbf{X}\mathbf{Y} - \mathbf{Z}\|_F^2 \quad \text{s.t. } Z_{ij} = M_{ij}, \quad \forall (i, j) \in \Omega$$

where $\mathbf{X} \in \mathbb{R}^{M \times L}$, $\mathbf{Y} \in \mathbb{R}^{L \times N}$, $\mathbf{Z} \in \mathbb{R}^{M \times N}$, and L is an estimate of min. rank.

- Advantage of adopting (\triangle) : When BCD is applied, each subproblem of (\triangle) has a closed-form solution:

$$\begin{aligned} \mathbf{X}^{r+1} &= \mathbf{Z}^r \mathbf{Y}^{rT} (\mathbf{Y}^r \mathbf{Y}^{rT})^\dagger, \\ \mathbf{Y}^{r+1} &= (\mathbf{X}^{r+1T} \mathbf{X}^{r+1})^\dagger (\mathbf{X}^{r+1T} \mathbf{Z}^r), \\ [\mathbf{Z}^{r+1}]_{i,j} &= \begin{cases} [\mathbf{X}^{r+1} \mathbf{Y}^{r+1}]_{i,j}, & \text{for } (i, j) \notin \Omega \\ M_{i,j}, & \text{for } (i, j) \in \Omega \end{cases} \end{aligned}$$

Applications — Maximizing A Convex Quadratic Function

- Consider maximizing a convex quadratic problem:

$$(\square) \quad \max_{\mathbf{x}} \frac{1}{2} \mathbf{x}^T \mathbf{Q} \mathbf{x} + \mathbf{c}^T \mathbf{x} \quad \text{s.t. } \mathbf{x} \in \mathcal{X}$$

where \mathcal{X} is a polyhedral set, and $\mathbf{Q} \succeq \mathbf{0}$.

- (\square) is equivalent to the following problem¹

$$(\triangle) \quad \max_{\mathbf{x}_1, \mathbf{x}_2} \frac{1}{2} \mathbf{x}_1^T \mathbf{Q} \mathbf{x}_2 + \frac{1}{2} \mathbf{c}^T \mathbf{x}_1 + \frac{1}{2} \mathbf{c}^T \mathbf{x}_2 \quad \text{s.t. } (\mathbf{x}_1, \mathbf{x}_2) \in \mathcal{X} \times \mathcal{X}$$

- When fixing either \mathbf{x}_1 or \mathbf{x}_2 , problem (\triangle) is an LP, thereby efficiently solvable.

¹The equivalence is in the following sense: If \mathbf{x}^* is an optimal solution of (\square) , then $(\mathbf{x}^*, \mathbf{x}^*)$ is optimal for (\triangle) ; Conversely, if $(\mathbf{x}_1^*, \mathbf{x}_2^*)$ is an optimal solution of (\triangle) , then both $\mathbf{x}_1^*, \mathbf{x}_2^*$ are optimal for (\square) .

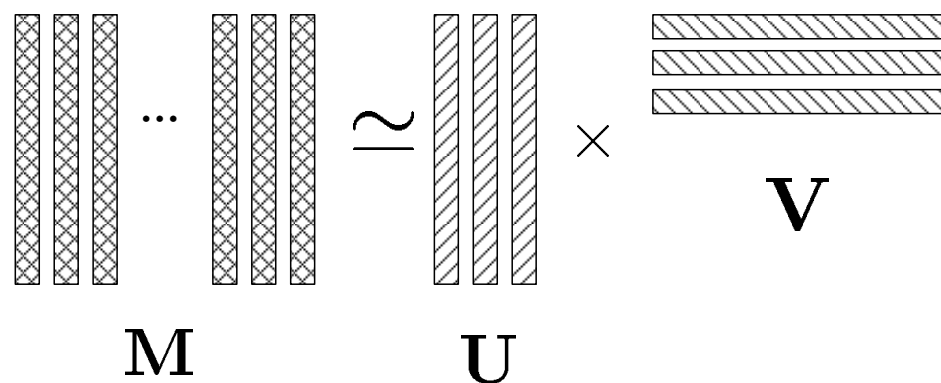
Applications — Nonnegative Matrix Factorization (NMF)

- NMF is concerned with the following problem [Lee-Seung]:

$$\min_{\mathbf{U} \in \mathbb{R}^{m \times k}, \mathbf{V} \in \mathbb{R}^{k \times n}} \|\mathbf{M} - \mathbf{UV}\|_F^2 \quad \text{s.t. } \mathbf{U} \geq \mathbf{0}, \mathbf{V} \geq \mathbf{0} \quad (15)$$

where $\mathbf{M} \geq \mathbf{0}$.

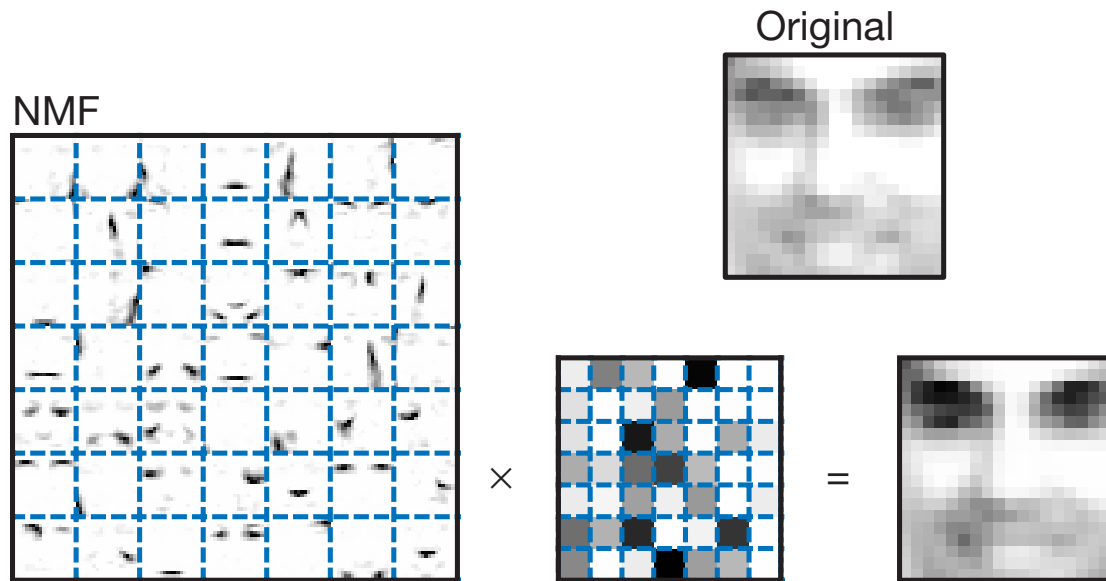
- Usually $k \ll \min(m, n)$ or $mk + nk \ll mn$, so NMF can be seen as a linear dimensionality reduction technique for nonnegative data.



NMF Examples

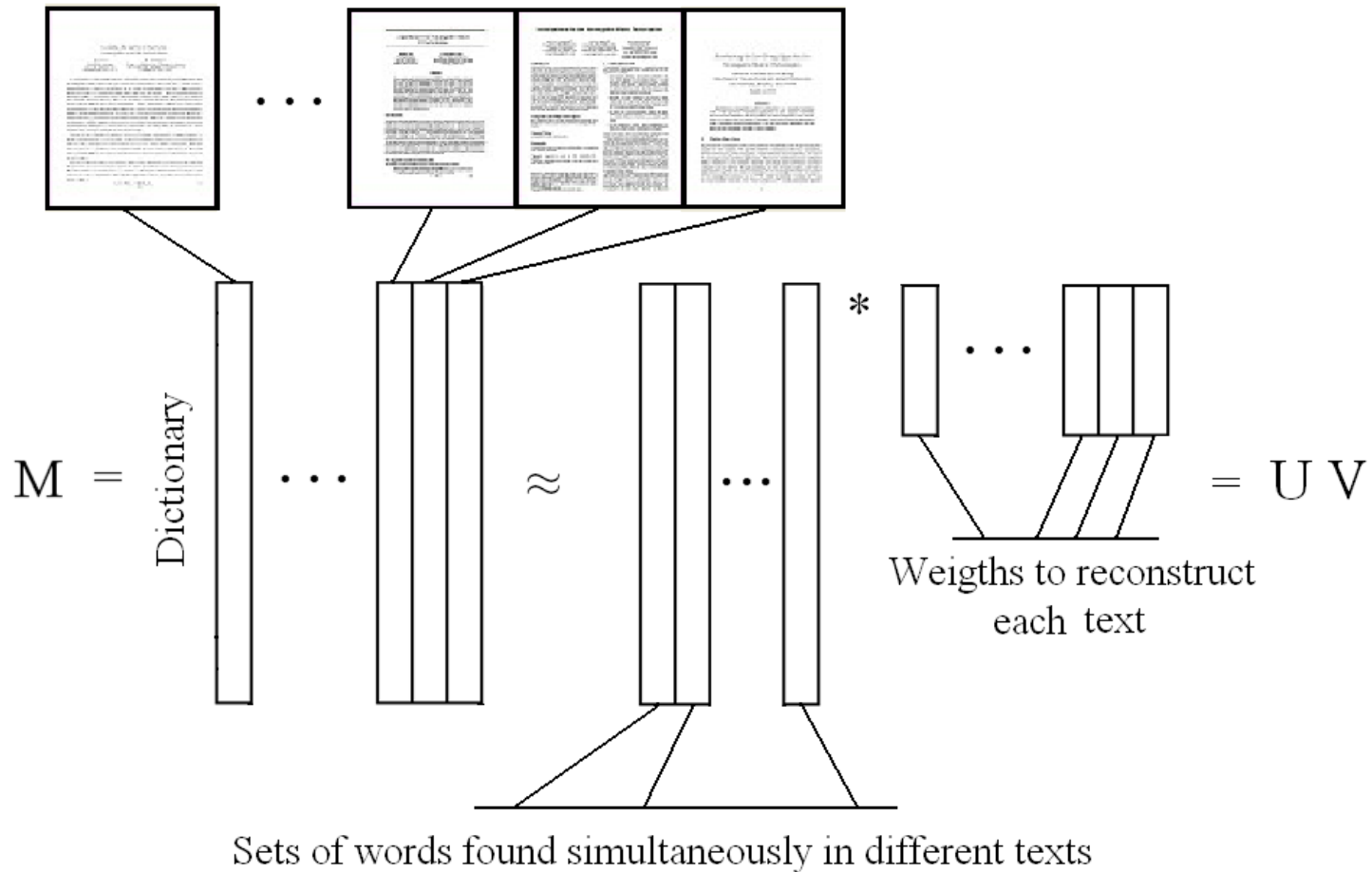
- **Image Processing:**

- $\mathbf{U} \geq \mathbf{0}$ constraints the basis elements to be nonnegative.
- $\mathbf{V} \geq \mathbf{0}$ imposes an additive reconstruction.



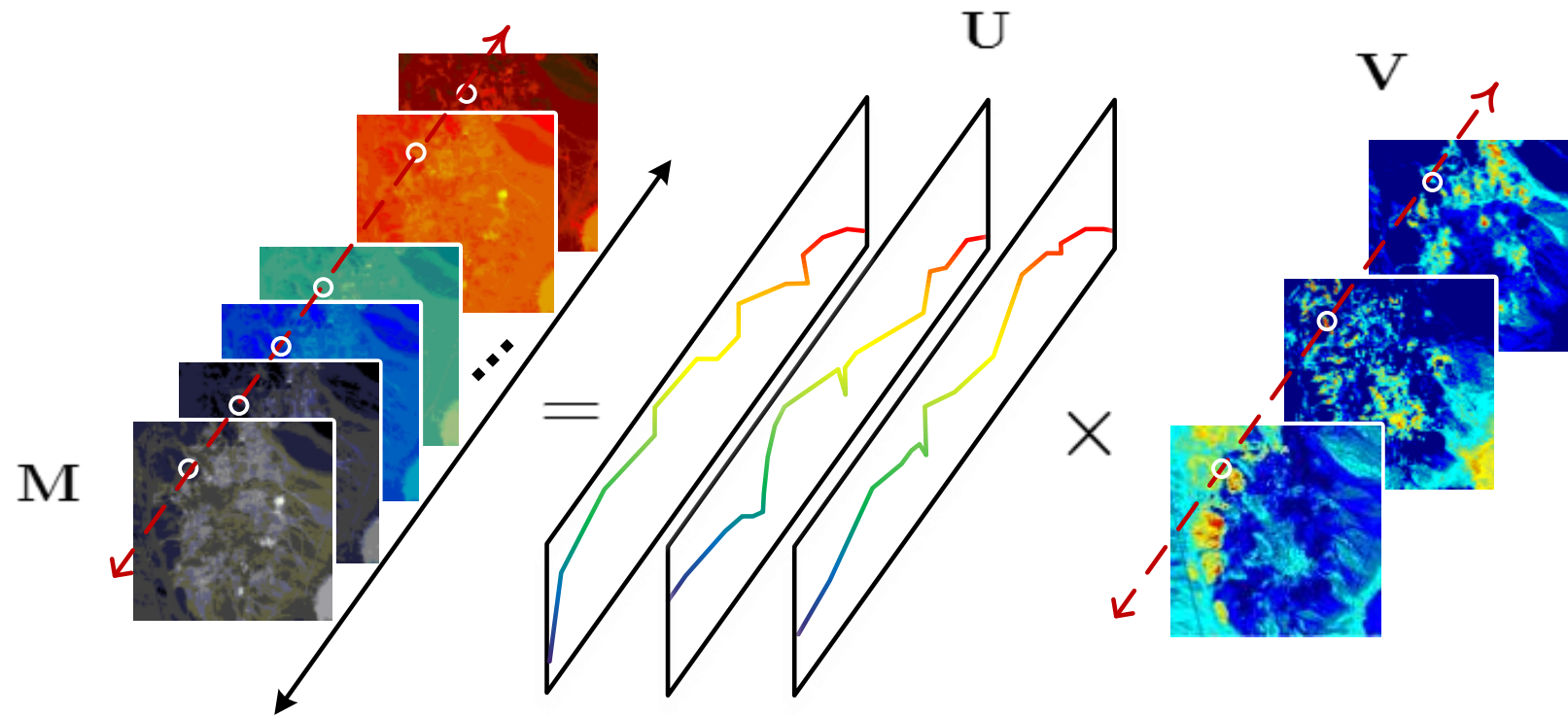
The basis elements extract facial features such as eyes, nose and lips.

• Text Mining



- Basis elements allow to recover different topics;
- Weights allow to assign each text to its corresponding topics.

• Hyperspectral Unmixing



- Basis elements U represent different materials;
- Weights V allow to know which pixel contains which material.

- Let's turn back to the NMF problem:

$$\min_{\mathbf{U} \in \mathbb{R}^{m \times k}, \mathbf{V} \in \mathbb{R}^{k \times n}} \|\mathbf{M} - \mathbf{UV}\|_F^2 \quad \text{s.t. } \mathbf{U} \geq \mathbf{0}, \mathbf{V} \geq \mathbf{0} \quad (16)$$

- Without “ $\geq \mathbf{0}$ ” constraints, the optimal \mathbf{U}^* and \mathbf{V}^* can be obtained by SVD.
- With “ $\geq \mathbf{0}$ ” constraints, problem (16) is generally NP-hard.
- When fixing \mathbf{U} (resp. \mathbf{V}), problem (16) is convex w.r.t. \mathbf{V} (resp. \mathbf{U}).
- For example, for a given \mathbf{U} , the i th column of \mathbf{V} is updated by solving the following NLS problem:

$$\min_{\mathbf{V}(:,i) \in \mathbb{R}^k} \|\mathbf{M}(:,i) - \mathbf{UV}(:,i)\|_2^2, \quad \text{s.t. } \mathbf{V}(:,i) \geq \mathbf{0}, \quad (17)$$

BCD Algorithm for NMF:

- 1: Initialize $\mathbf{U} = \mathbf{U}^0$, $\mathbf{V} = \mathbf{V}^0$ and $r = 0$;
- 2: **repeat**
- 3: solve the NLS problem

$$\mathbf{V}^* \in \arg \min_{\mathbf{V} \in \mathbb{R}^{k \times n}} \|\mathbf{M} - \mathbf{U}^r \mathbf{V}\|_F^2, \quad \text{s.t. } \mathbf{V} \geq \mathbf{0}$$

- 4: $\mathbf{V}^{r+1} = \mathbf{V}^*$;
- 5: solve the NLS problem

$$\mathbf{U}^* \in \arg \min_{\mathbf{U} \in \mathbb{R}^{m \times k}} \|\mathbf{M} - \mathbf{U} \mathbf{V}^{r+1}\|_F^2, \quad \text{s.t. } \mathbf{U} \geq \mathbf{0}$$

- 6: $\mathbf{U}^{r+1} = \mathbf{U}^*$;
- 7: $r = r + 1$;
- 8: **until** some convergence criterion is met

Outline

- Majorization Minimization (MM)
 - Convergence
 - Applications
- Block Coordinate Descent (BCD)
 - Applications
 - Convergence
- Summary

BCD Convergence

- The idea of BCD is to divide and conquer. However, there is no free lunch; BCD may get stuck or converge to some point of no interest.

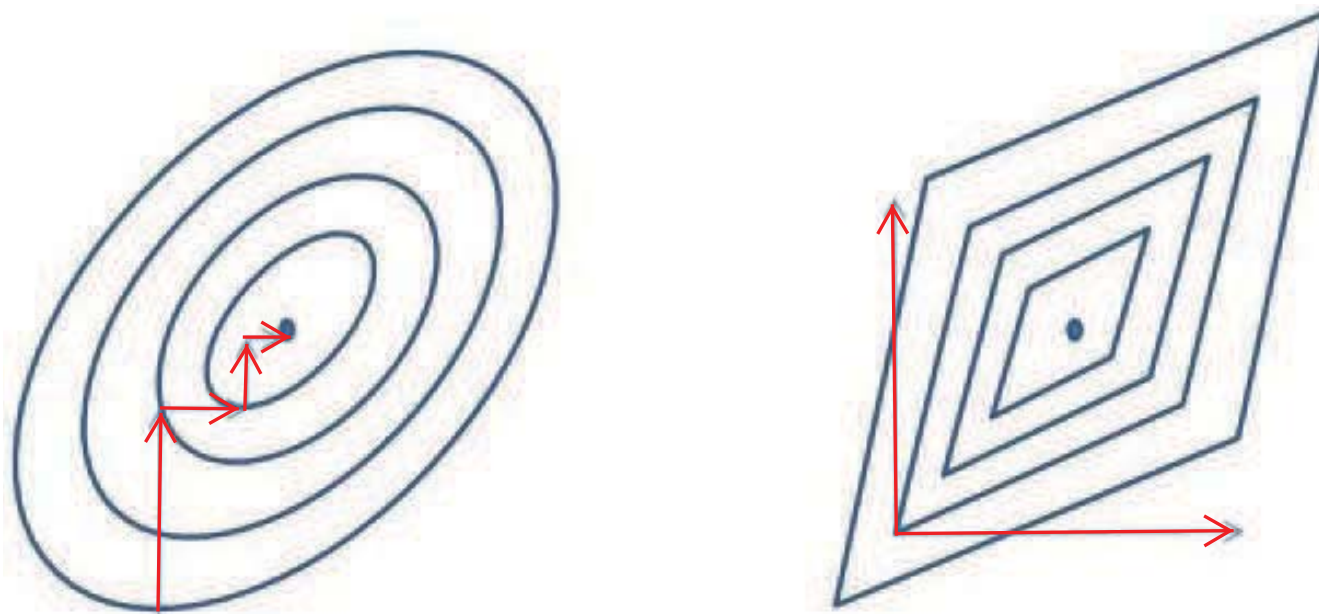


Figure 6: BCD for smooth/non-smooth minimization.

BCD Convergence

$$\min_{\mathbf{x}} f(\mathbf{x}) \quad \text{s.t. } \mathbf{x} \in \mathcal{X} = \mathcal{X}_1 \times \mathcal{X}_2 \times \dots \times \mathcal{X}_m \subseteq \mathbb{R}^n \quad (18)$$

- A well-known BCD convergence result due to Bertsekas:

Theorem 2. (*[Bertsekas]*) Suppose that f is continuously differentiable over the convex closed set \mathcal{X} . Furthermore, suppose that for each i

$$g_i(\boldsymbol{\xi}) \triangleq f(\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_{i-1}, \boldsymbol{\xi}, \mathbf{x}_{i+1}, \dots, \mathbf{x}_m)$$

is *strictly convex*. Let $\{\mathbf{x}^r\}$ be the sequence generated by BCD method. Then every limit point of $\{\mathbf{x}^r\}$ is a stationary point of problem (18).

- If \mathcal{X} is (convex) compact, i.e., closed and bounded, then strict convexity of $g_i(\boldsymbol{\xi})$ can be relaxed to having a unique optimal solution.

- Application: Iterative water-filling for MIMO MAC sum capacity max.:

$$(\Delta) \quad \max_{\{\mathbf{Q}_k\}_{k=1}^K} \log \det \left(\sum_{k=1}^K \mathbf{H}_k \mathbf{Q}_k \mathbf{H}_k^H + \mathbf{I} \right), \quad \text{s.t. } \text{Tr}(\mathbf{Q}_k) \leq P_k, \quad \mathbf{Q}_k \succeq \mathbf{0}, \quad \forall k$$

- Iterative water-filling converges to a global optimal solution of (Δ) , because
 - BCD subproblem is strictly convex (assuming full column rankness of \mathbf{H}_k);
 - \mathcal{X}_k is a convex closed subset;
 - (Δ) is a convex problem, so stationary point \implies global optimal solution

Generalization of Bertsekas' Convergence Result

- Generalization 1: Relax Strict Convexity to Strict Quasiconvexity² [Grippi-Sciandrone]

Theorem 3. *Suppose that the function f is continuously differentiable and **strictly quasiconvex** with respect to x_i on \mathcal{X} , for each $i = 1, \dots, m - 2$ and that the sequence $\{x^r\}$ generated by the BCD method has limit points. Then, every limit point is a stationary point of problem (18).*

- Application: Low-Rank Matrix Completion

$$(\Delta) \quad \min_{\mathbf{X}, \mathbf{Y}, \mathbf{Z}} \frac{1}{2} \|\mathbf{X}\mathbf{Y} - \mathbf{Z}\|_F^2 \quad \text{s.t.} \quad Z_{ij} = M_{ij}, \quad \forall (i, j) \in \Omega$$

- $m = 3$ and (Δ) is strictly convex w.r.t. $\mathbf{Z} \implies$ BCD converges to a stationary point.

² f is strictly quasiconvex w.r.t. $x_i \in \mathcal{X}_i$ on \mathcal{X} if for every $\mathbf{x} \in \mathcal{X}$ and $\mathbf{y}_i \in \mathcal{X}_i$ with $\mathbf{y}_i \neq \mathbf{x}_i$ we have

$$f(\mathbf{x}_1, \dots, t\mathbf{x}_i + (1 - t)\mathbf{y}_i, \dots, \mathbf{x}_m) < \max \{f(\mathbf{x}), f(\mathbf{x}_1, \dots, \mathbf{y}_i, \dots, \mathbf{x}_m)\}, \quad \forall t \in (0, 1).$$

- Generalization 2: Without Solution Uniqueness

Theorem 4. Suppose that f is *pseudoconvex*³ on \mathcal{X} and that $\mathcal{L}_\mathcal{X}^0 := \{\mathbf{x} \in \mathcal{X} : f(\mathbf{x}) \leq f(\mathbf{x}^0)\}$ is compact. Then, the sequence generated by BCD method has limit points and every limit point is a *global minimizer* of f .

- Application: Iterative water-filling for MIMO-MAC sum capacity max.

$$\begin{aligned} \max_{\{\mathbf{Q}_k\}_{k=1}^K} \quad & \log \det \left(\sum_{k=1}^K \mathbf{H}_k \mathbf{Q}_k \mathbf{H}_k^H + \mathbf{I} \right) \\ \text{s.t.} \quad & \text{Tr}(\mathbf{Q}_k) \leq P_k, \quad \mathbf{Q}_k \succeq \mathbf{0}, \quad k = 1, \dots, K \end{aligned}$$

- f is convex, thus pseudoconvex;
- $\{\mathbf{Q}_k \mid \text{Tr}(\mathbf{Q}_k) \leq P_k, \mathbf{Q}_k \succeq \mathbf{0}\}$ is compact;
- iterative water-filling converges to a globally optimal solution.

³ f is pseudoconvex if for all $\mathbf{x}, \mathbf{y} \in \mathcal{X}$ such that $\nabla f(\mathbf{x})^T(\mathbf{y} - \mathbf{x}) \geq 0$, we have $f(\mathbf{y}) \geq f(\mathbf{x})$. Notice that “convex \subset pseudoconvex \subset quasiconvex”.

- Generalization 3: Without Solution Uniqueness, Pseudoconvexity and Compactness

Theorem 5. *Suppose that f is continuously differentiable, and that \mathcal{X} is convex and closed. Moreover, if there are only **two** blocks, i.e., $m = 2$, then every limit point generated by BCD is a stationary point of f .*

- Application: NMF

$$\min_{\mathbf{U} \in \mathbb{R}^{m \times k}, \mathbf{V} \in \mathbb{R}^{k \times n}} \|\mathbf{M} - \mathbf{UV}\|_F^2 \quad \text{s.t. } \mathbf{U} \geq \mathbf{0}, \mathbf{V} \geq \mathbf{0}$$

- Alternating NLS converges to a stationary point of the NMF problem, since
 - the objective is continuously differentiable;
 - the feasible set is convex and closed;
 - $m = 2$.

Summary

- MM and BCD have great potential in handling nonconvex problems and realizing fast/distributed implementations for large-scale convex problems;
- Many well-known algorithms can be interpreted as special cases of MM and BCD;
- Under some conditions, convergence to stationary point can be guaranteed by MM and BCD.

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